Quantifying the Utility of Noisy Reviews in Stopping Information Cascades

Tho Ngoc Le¹, Vijay G. Subramanian², Randall A. Berry¹

Abstract—In models of social learning where rational agents can observe other agents' actions, information cascades are said to occur when agents ignore their own private information and blindly follow the actions of other agents. It is well known that in some cases, incorrect cascades happen with positive probability leading to a loss in social welfare. Having agents provide reviews in addition to their actions provides one possible way to avoid such "bad cascades." In this paper, we study one such model where agents sequentially decide whether or not to purchase a good, whose true value is either good or bad. If they purchase the good, agents also leave a review, which may be noisy. Conditioning on the underlying state of the world, we study the impact of such reviews on the asymptotic properties of cascades. For a good underlying state, we propose an algorithm that utilizes number theory and Markov analysis to solve for the probability of wrong cascade. We discover that depending on the review quality, reviews may change the probability of a wrong cascade in a non-monotonic manner. On the other hand, for a bad underlying state, the agents always eventually reach a correct cascade; we use a martingale analysis to bound the time until this happens.

I. Introduction

On-line platforms provide an easy way for people to attempt to learn from others before making a new decision. Such "social learning" has long been studied by economists as a game among Bayesian agents. In the simplest setting, these agents sequentially make a binary decision based on their own beliefs, which are in turn a function of their own private information as well as observations of the decisions of previous agents. A key result, first shown in [2] and [3], is that such models exhibit *herding* or *information cascades*. This refers to a case where some agents ignore their private information and follow the actions of the previous agents. Moreover, for the models in [2] and [3], once a cascade starts, all subsequent agents also cascade. Though individually optimal, this may result in the agents making a choice that is not socially optimal, which we refer to as a "wrong cascade."

A wrong cascade occurs because agents observe the actions of other agents *before* the other agents receive their pay-offs, and so these actions reflect the agents' estimates of the true pay-off and not the true pay-off itself. Indeed, if agents instead were able to see the true pay-off obtained by others, then as shown in [9] there would never be an incorrect cascade in which agents buy a bad product. The use of

reviews and on-line recommendation systems can be viewed as an attempt to provide other agents with this information. However, due for example to user errors, such reviews may only be a noisy representation of this information (instead of the true pay-off as in [9]).

The goal of this paper is to study social learning in the presence of such noisy reviews. More precisely, we consider a variation of the models in [2], [3], where agents have the option to either buy or not buy a given item, whose true value is one of two binary states (good or bad). In addition to the actions of the previous agents, agents also see a history of reviews before making their decisions. However, these reviews are not a perfect indication of the true state of the good due to two effects: first, as we have already mentioned, these reviews are noisy, and second, agents can only leave a review if they buy the good and so no additional information is given for agents that choose not to buy.¹

Adding reviews is a way of changing the information structure in [2], [3]. A number of other ways have been considered that also change this structure such as changing the underlying network structure among the agents, e.g. [8], or changing the signal structure, e.g. [6]. In prior work ([10], [11]), we considered a variation of the information structure, where agent observed noisy observations of the actions of others. This led to the following counter-intuitive result: the probability of incorrect herding is non-monotonic in the noise level. In other words, in some cases, more noise is actually beneficial. In this paper, we again seek to study how variations in the noise level affect the agents' behaviors. However, here, agents perfectly observe the actions of previous agents and the only noise is in the reviews. Additionally, since only agents who buy the good can submit reviews, this leads to an asymmetry in the model that was not present in [10], [11]).

We presented an initial analysis of this model in [12]. There it was shown that the asymmetry in reviewing leads to an asymmetry in the resulting users' behaviors depending on the underlying state of the product being either good or bad. Here we present a more refined analysis of these two cases. Conditioned on the state of the product being good, we study the probability of an incorrect cascade. We give an algorithm based on a number theoretic arguments that enables characterizing this probability for a much larger set of parameter settings. Using this we then can characterize the behavior of the wrong cascade probability as a function of the noise level and show that this is a highly non-monotonic

¹Tho Ngoc Le and Randall A. Berry is with the Department of Electrical Engineering and Computer Science, Northwestern University, Evanston, IL 60208, USA. thole2012@u.northwestern.edu, rberry@eecs.northwestern.edu

²Vijay G. Subramanian is with the Department of Electrical Engineering and Computer Science, University of Michigan, Ann Arbor, MI 48109, USA vgsubram@umich.edu

¹For example, many on-line platforms such as Amazon.com indicate verified purchase reviews; in our mode only such reviews are considered.

and discontinuous function, so that in some cases decreasing the reviews noise can lead to a higher probability of a wrong cascade. Conditioned on the state being bad, we instead focus on the expected time until a correct cascade occurs. Using martingale techniques, we give bounds on the expected time until a correct cascade happens. We compare these bounds with simulations and offer an algorithm to improve the lower bound numerically.

Another strand of related work is the literature on "word-of-mouth" learning (e.g. [4], [5], [7]) in which agents can communicate information about payoff of past actions. However, these models consider different settings (e.g. naive rule-of-thumb decision-based, random sampling of population); while our paper assumes that fully-rational agents can observe all past actions and reviews.

We organize this paper as follows. In Section II we specify our model. The main results are presented in sections III and IV for the cases where the value of product is "good" and "bad," respectively. We conclude in Section V.

II. MODEL

We consider a model similar to [12] in which there is a countable population of agents, indexed $n=1,2,\ldots$ with the index reflecting the time and the order in which agents act (given exogenously). There is a new product (item) with a true value (V) that can be either good (G) or bad (B); both possibilities are assumed to be equally likely and the value is the same for all agents. Each agent n has a one-time action choice A_n of saying either "Yes" (Y) or "No" (N) to this item. Assume each agent n has prior knowledge about the true value V via a private signal $S_n \in \{1 \text{ (high)}, 0 \text{ (low)}\}$. For each agent n who chooses $A_n = Y$ submits a review $R_n \in \{G \text{ (Good)}, B \text{ (Bad)}\}$ representing his experience with the item after purchasing. On the other hand, if agent n chooses $A_n = N$, he does not submit a review.

We consider a homogeneous population where, conditioned on V, the private signals and reviews are i.i.d. across all agents. Assume the probability that a private signal (resp. a review) aligns with V is $p \in (0.5,1)$ (resp. $\delta \in [0.5,1]$), i.e., the distributions of the signals and reviews are given as:

$$\begin{split} \mathbb{P}[S_n = 1 | V = B] &= \mathbb{P}[S_n = 0 | V = G] = 1 - p, \\ \mathbb{P}[S_n = 1 | V = G] &= \mathbb{P}[S_n = 0 | V = B] = p, \text{ and if } A_n = Y, \\ \mathbb{P}[R_n = G | V = G] &= \mathbb{P}[R_n = B | V = B] = \delta, \\ \mathbb{P}[R_n = G | V = B] &= \mathbb{P}[R_n = B | V = G] = 1 - \delta. \end{split}$$

Since $p \in (0.5, 1)$, the private signals are informative, but not revealing; we call p the signal quality. On the other hand, δ denotes the review's strength. The review and the private signal are assumed to be conditionally independent given V. Let $\mathcal{R}_n = R_n$ when $A_n = Y$ and $\mathcal{R}_n = *$ when $A_n = N$. The history after agent n decides is written as $H_n = \{A_1, \mathcal{R}_1, \ldots, A_n, \mathcal{R}_n\}$; we assume that H_n is public information to subsequent agents. The agents

are Bayes-rational whose decisions are based on their own private signals and public information. Each agent n updates his posterior belief about the true value V using his private signal S_n , the actions A_1, \ldots, A_{n-1} , and the reviews $\mathcal{R}_1, \ldots, \mathcal{R}_{n-1}$.

A. Public likelihood ratio as a Markov process

Let q=1-p. Agents' decisions are based on Bayes' updates of the posterior probability of V=B versus V=G given the observed history H_n . However, due to the independence of signals from the public history, agent n+1 can instead compare the public likelihood ratio, ℓ_n , and his private belief β_{n+1} , of V=B versus V=G. Since V being $\{B,G\}$ equally likely, $\ell_0=1$ and we can rewrite ℓ_n in its alternate form:

$$\ell_n = \frac{\mathbb{P}[H_n | V = B]}{\mathbb{P}[H_n | V = G]}, \text{ and } \beta_{n+1} = \frac{\mathbb{P}[S_{n+1} | V = B]}{\mathbb{P}[S_{n+1} | V = G]}. \quad (1)$$

The higher ℓ_n is, the more likely that V=B. Moreover, since H_n is public information, ℓ_n can be updated as:

• If agent n follows his own signal then:

$$\ell_n = \begin{cases} \frac{p}{q} \ell_{n-1}, & \text{if } A_n = N, \\ \frac{q}{p} \frac{1-\delta}{\delta} \ell_{n-1}, & \text{if } A_n = Y, \mathcal{R}_n = G, \\ \frac{q}{p} \frac{\delta}{1-\delta} \ell_{n-1}, & \text{if } A_n = Y, \mathcal{R}_n = B. \end{cases}$$
 (2)

• Otherwise, if agent n cascades then:

$$\ell_n = \begin{cases} \ell_{n-1}, & \text{if } A_n = N, \\ \frac{1-\delta}{\delta}\ell_{n-1}, & \text{if } A_n = Y, \mathcal{R}_n = G, \\ \frac{\delta}{1-\delta}\ell_{n-1}, & \text{if } A_n = Y, \mathcal{R}_n = B. \end{cases}$$
(3)

Moreover, as shown in the next lemma, given ℓ_n one can determine if an agent cascades or not. Thus, $\{\ell_n\}$ is a Markov process. Moreover, this is also true if in addition we condition on each value of V^4 . On the other hand, $\beta_{n+1}=q/p$ (resp. p/q) if $S_{n+1}=1$ (resp. $S_{n+1}=0$).

B. Agents' decision rule and cascades' condition

Let a_n and r_n be two integer random variables denoting the two differences in actions (#Y - #N) and reviews (#G - #B), respectively. Note that while a_n excludes the actions caused by both types of cascades (since then the cascading actions provide no information), r_n is unchanged whenever a review is not made due to an agent not buying.

Lemma 1. Define $x = \log_{\frac{\delta}{1-\delta}}(p/q) \in (0,\infty)$ for $\delta \in (0.5,1)$. Then:

- 1) $\ell_n = (q/p)^{h_n}$, where the exponent $h_n = a_n + \frac{1}{r}r_n$,
- 2) Conditioned on V, (a_n, r_n) and h_n are 2-D and 1-D Markov chains, respectively, for $n \ge 0$,
- 2) Agent n+1 cascades Y if $h_n > 1$, cascades N if $h_n < -1$, and follows his signal if $h_n \in [-1, 1]$.

Proof. 1) By (2) and (3), $\ell_n = (q/p)^{a_n} ((1-\delta)/\delta)^{r_n}$, thus h_n can be written in terms of a_n and r_n as above.

²The motivation being, while signal quality reflects a product's marketing efficiency, the review strength is a consequence of product reliability, e.g., due to manufacturing.

³For simplicity, we assume indifferent agents follow their own signals. ⁴This is an extension of results from [6].

- 2) This is a direct consequence of the fact that $\{\ell_n\}$ is a Markov process and that, from the first property, there is a 1-1 correspondence between ℓ_n and h_n . Further, since a_n and r_n are integer-valued it follows that h_n only takes on a countable number of values. Without reviews a similar Markov chain was used in [10].
- 3) Since agent n+1 makes his decision by comparing $\ell_n\beta_{n+1}$ to 1, agent n+1 cascades Y if $\ell_n < q/p$, cascades N if $\ell_n > p/q$, and follows his signal if $\ell_n \in [q/p, p/q]$. By 1), this is translated to the given condition on h_n .

Note that x is an indicator of how strong the reviews are with respect to the signals. That is, the lower x is, the stronger the reviews are relative to the signals. For a generic x, the dynamics of the process $\{\ell_n\}$ can be studied by investigating the 2-D Markov chain (a_n, r_n) . However, for special values of x, this can be simplified to certain extents. We will study two of such scenarios in Section III.

C. Asymmetry by different types of cascade and item quality

This model exhibits asymmetric behaviors with respect both to the types of cascades (Y and N), and to the true value V of the item. In particular, the arrival of new information (reviews) depends on the action chosen by each agent. We first highlight a key difference between Y and N cascades in the following two properties.

Property 1. Once a Y cascade starts, there is a positive probability that it ends (unless the review are perfect).

If agent n faces $h_{n-1} > 1$, he chooses $A_n = Y$ regardless of his signal, and thus initiates a Y cascade. Such a cascade can end if subsequent agents submit a sufficient number of bad reviews, e.g., if $\mathcal{R}_n = B$, then $h_n = h_{n-1} - \frac{1}{x}$ could be below 1, which induces agent n+1 to use his signal. Furthermore, if x is sufficiently small then agent n's bad review can make $h_n < -1$, so that agent n+1 starts a N cascade. The dynamics of a Y cascade, once it gets started, are determined solely by the reviews process (and it does not depend on the signals). Regardless of the time a Y cascade was initiated, it can be broken by a sufficiently long sequence of bad reviews. Thus, the history process $\{H_n\}$ could include sample paths where Y cascades start and stop multiple times.

On the other hand, once $h_n < -1$, a N cascade starts, it lasts forever. This is because agents who choose N do not generate reviews; thus, the likelihood ratio stays constant as soon as any agent cascades to N. Subsequent agents are left in the same state as the one who initiated the cascade; thus make the same action choice. We summarize this in the following Property 2.

Property 2. Once a N cascade happens, it lasts forever.

Next we give two properties that show the differences between a good and a bad product.

Property 3. For V = G, a wrong cascade happens with positive probability

This comes as a result of the existence of the absorbing states for a wrong cascade. For example, if the first two agents have low signals, they both choose N; therefore no review is collected. As a result, all subsequent agents are drawn into a N cascade, which is irreversible. This possibility cannot be avoided by adjusting the reviews strength, δ , even to perfect quality. In case the reviews are perfect, we would still need a non-cascading agent who has a H signal for his review to be submitted.

Though a wrong cascade is possible, for V=G, it is more likely that there will be an abundance of information, since each agent that chooses Y also creates a new review. Since reviews are independent of signal, when V=G more agents choose Y, and new information begets further new information. In other words, when V=G the underlying Markov process has a drift toward the correct cascade, but there is no absorbing state on that side since h_n is unbounded above. However, multiple absorbing states for wrong cascade might exist. For V=G, the quantity of interest is the probability of wrong (N) cascade, which is a function of both p and δ . We will discuss this scenario in section III.

On the other hand, when V=B, this model exhibits a different set of behaviors. As more agents purchase the item, more and more reviews are collected. Since reviews are informative, subsequent agents can track the difference in the number of reviews to learn the true value of V eventually. In other words, while there are only trapping states for correct cascade, the drift also leans toward this side. We summarize this result below.

Property 4. For V = B and $\delta > 0.5$, a wrong cascade can never happen.¹

Thus, for V=B correct cascade happens with probability 1. In this scenario, we are interested in the distribution of the time (i.e. the number of agents) until a correct cascade happens. This will be studied in section IV.

III. Probability of wrong cascade for $V={\cal G}$

In previous section, we discussed that wrong (N) cascades could happen if the product is good. In this section, we determine the probability of this happening. For a fixed p, as x varies the conditions on a_n and r_n when cascades happen also change. As a result, the underlying Markov chains have different structures (both in terms of the state spaces and the transition probabilities). Despite the complexity of these dynamics for a generic x, interesting and non-intuitive insights can be drawn by looking at special values of x. In one example, for any rational x, many states of (a_n, r_n) can be mapped to one single state of (h_n) ; thus it is sufficient to study the reduced 1-D Markov chain (h_n) . This is generally not possible for any real value of x since there would be a 1-1 mapping between the states in (a_n, r_n) and (h_n) ; thus this prevents the simplification of the state space. However, in another example when x is real and x < 1/3, the state space of (a_n, r_n) can also be simplified to obtain analytical results. In particular, we consider two scenarios that facilitate simplification of the underlying state space of (a_n, r_n) :

¹Note if $\delta = 0.5$, then reviews are useless, in which case wrong cascades can occur as in [2].

- 1) x is a rational number in $(0, \infty)$, and
- 2) x is any real number in (0, 1/3].

A. x = i/j for positive integers i, j and gcd(i, j) = 1

From the discussions at the beginning of this section, it is sufficient in this case to consider the 1-D Markov chain (h_n) . Let \mathbb{P}_s be the asymptotic probability of wrong cascade starting from the state $h_0=s$. We want to find \mathbb{P}_0 . Given i, consider the finite set $\mathscr{A} \stackrel{\triangle}{=} \{-1, -\frac{i-1}{i}, \dots, \frac{i-1}{i}, 1\}$. It is obvious that \mathscr{A} is the set of all possible values that $h_n=a_n+\frac{j}{i}r_n$ can take in [-1,1]. Depending on the value of x, the following Lemma 2 further reduces the set of accessible states for $h_n \in [-1,1]$ to different subsets of \mathscr{A} .

Lemma 2. Assume x = i/j is rational, where i, j are positive integers with gcd(i, j) = 1:

1) If
$$x \le 1/3$$
, or if $x \in \{1/2, 1\}$, $h_n \in \{-1, 0, 1\}$,
2) If $1/3 < x < 1/2$, let $z = j \mod i$ and $k = \lfloor i/z \rfloor$ Then $h_n \in \{-1, 0, 1, -\frac{z}{i}, -\frac{2z}{i}, \dots, -\frac{kz}{i}, \frac{i-z}{i}, \frac{i-2z}{i}, \dots, \frac{i-kz}{i}\}$,
3) If $x > 1/2$, h_n takes all the values in \mathscr{A} .

Proof Idea. The proof uses number theoretic arguments to find what values in \mathscr{A} can be obtained (see Appendix). \square

As a consequence of Lemma 2, we can numerically solve for \mathbb{P}_0 . The idea is based on Markov analysis where one can write down a system of of linear equations (LEs) with the set of variables being \mathbb{P}_{h_n} for all accessible states h_n . Since there is no absorbing state for a Y cascade, h_n is not upperbounded and the accessible state space is infinite. However, once $h_n > 1$ the state transitions dynamics are simplified to a birth-death process; thus any variable \mathbb{P}_{h_n} where $h_n > 1$ can be expressed in terms of the corresponding variables where $h_n \in \mathscr{A}$. Therefore, the number of equations is finite (at most 2i+1). We propose the following Algorithm 1 to construct the system of linear equations and solve for \mathbb{P}_0 :

Algorithm 1 Wrong cascade probability at rational x

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Input: V=1,\ p,\ x=i/j,\ gcd(i,j)=1
Output: LEs and solution \mathbb{P}_0
\delta\leftarrow 1/\left(1+(q/p)^{1/x}\right),\ q\leftarrow 1-p,\ \alpha\leftarrow (1-\delta)/\delta
Initialize \mathscr{A}\leftarrow \{-1,-(i-1)/i,\ldots,(i-1)/i,1\}
\mathscr{A}\supseteq\mathscr{A}'\leftarrow \text{ accessible states in } [-1,1] \text{ (Lemma 2).}
for h_n=s\in\mathscr{A}' do
s_L\leftarrow s-1,s_{HB}\leftarrow s+1-j/i,s_{HG}\leftarrow s+1+j/i
c_1\leftarrow \text{ min number of steps from } s_{HG} \text{ to } s_1\in\mathscr{A}'.
if s_{HB}>1 then
c_2\leftarrow \text{ min number of steps from } s_{HB} \text{ to } s_2\in\mathscr{A}'
Eq_s\leftarrow \mathbb{P}_s=q\mathbb{P}_{s_L}+p\delta\alpha^{c_1}\mathbb{P}_{s_1}+p(1-\delta)\alpha^{c_2}\mathbb{P}_{s_2}
Add equation Eq_s to the system of LEs
Solve for \mathbb{P}_0 and return.
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Using Algorithm 1, numerical values for the probability of wrong cascade can be solved for, given x rational. Fig. 1 compares those numerical values to simulation results.

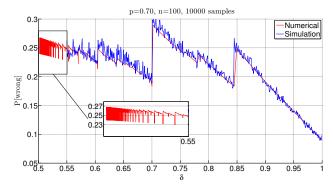


Figure 1: Wrong cascade probability for V = 1.

Both numerical and simulation results in Fig. 1 show that the probability of wrong cascade is not monotonic in the review quality δ . As δ varies in [0.5,1], there are points of discontinuities resulting from the changes in the state space and the transition probabilities of the underlying Markov chains (h_n) For low review quality, simulation fails to evaluate \mathbb{P}_0 since the mean number of agents, n, needed until a N cascade happens approaches infinity as $\delta \to 0.5$.

As a consequence of Lemma 2, for certain values of x the state space is simplified enough and we can obtain the closed-form expressions for the wrong cascade probability. In particular, Proposition 1 in [12] showed that for x=1 and x=1/2, $\mathbb{P}_0=(q/p)^2$. Moreover, when there are no reviews, a result from [2] gives $\mathbb{P}_0=(q/p)^2/\left[(q/p)^2+1\right]<(q/p)^2$. Thus, having reviews with strength equal or double the signal quality strictly increases the probability of wrong cascades. Further from Fig. 1, this is true for any $1/2 \le x \le 1$.

B. x is any real value in (0, 1/3].

In this section, we present the second scenario when the state space (a_n, r_n) can also be simplified. In particular, we look at the cases when reviews are at least three times stronger than the private signals. For any real value of x in this region, the state transitions of the underlying 2-D Markov chains are shown in Fig. 2, where the first and second coordinates denote r_n and a_n , respectively. By Proposition 2 in [12], we can obtain the closed-form expression for the probability of wrong cascade:

$$\mathbb{P}_0 = [1 - p(2\delta - 2p\delta + 2p - p/\delta)] / [1 - 2pq(1 - \delta)]$$
 (4)

which is decreasing in δ . This is illustrated in Fig. 3 for x=1/5,1/10. For all values of x in this figure, the probability of wrong cascade decreases in the signal quality p. Moreover, except for reviews with perfect accuracy, one would prefer having no reviews for enough low signal quality.

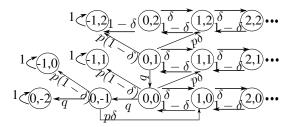


Figure 2: States transitions for V = G, and x < 1/3.

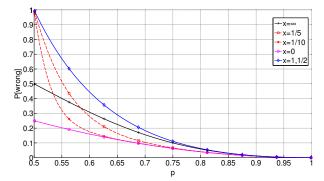


Figure 3: Wrong cascade probability for V = G.

The above conclusions can be explained by the discontinuity of the slopes for different curves in the above figure as $p \to 0.5$. With perfect reviews, the slope as $p \to 0.5$ is -1. With no review, the corresponding slope is at -2. However, as long as the platform generates reviews with strength δ bounded away from 0.5, the probabilities of wrong cascade follow the set of dashed curves shown, with slopes bounded away from -2. In particular, these slopes can be studied using (4) by setting $p=0.5+\epsilon$ and $x=C\epsilon$ where $\epsilon \to 0, C>0$. When C is fixed, δ is bounded away from 0.5; this yields a slope of $-\infty$ as $\epsilon \to 0$. When $C \to \infty$ and x < 1/3, $\delta \to 0.5$; and the slopes vary in $(-\infty, -8)$. Finally, if $C \to 0$, $\delta \to 1$; the slopes approach -1, which is exactly the slope of the perfect reviews scenario.

IV. Time until correct cascade for V=B

In section II, we argued that for a bad product, only a correct (N) cascade can happen, so that it lasts forever once it occurs. In this section, we examine both the upper and lower bounds on the expected time until correct cascades. In the following let $n \geq 0$. Conditioned on $V \in \{G, B\}$, let $\{\mathscr{F}_n^V\}$ be the sequence of σ -algebras generated by $\{H_n\}$. Similar to [6] and [8] where reviews do not exist, in our model the Markov process $\{\ell_n\}$ also exhibits the martingale property. In Section IV of [12] we showed that $\{1/\ell_n\}$ (resp. $\{\ell_n\}$) is a martingale process conditioned on V = B (resp. V = G) adapted to the filtration $\{\mathscr{F}_n^B\}$ (resp. $\{\mathscr{F}_n^G\}$). Moreover, let X and Y be two random variables representing the increments $\Delta h_n = h_{n+1} - h_n$ for h_n in [-1, 1] and $h_n > 1$, respectively. Let $f_1(\lambda)$ and $f_2(\lambda)$ be their corresponding moment generating functions (MGFs), where λ is a real variable. Let $\rho=\max(f_1(\lambda),f_2(\lambda))$ and define the random process $\{M_n\}=\left\{\frac{\mathrm{e}^{\lambda h_n}}{\rho^n}\right\}$. Using techniques from [1], in [12] we showed that $\{M_n\}$ is a super-martingale adapted to $\{\mathscr{F}_n^B\}$. Let $\tau=\min\{n\geq 0:h_n<-1\}$ be the stopping time when a N cascade happens. Now we use these results to bound the expected time until correct cascade, $\mathbb{E}[\tau]$.

A. Upper bound on $\mathbb{E}[\tau]$

Proposition 1. $\mathbb{E}[\tau] \leq e^{\lambda}/(1-\rho)$, where $0 < \rho < 1$, $\lambda \in (0, \ln(p/(1-p))$.

Proof. From Proposition 3 in [12], the tail distribution is upper-bounded by: $\mathbb{P}[\tau > n] \leq e^{\lambda} \rho^n$. For feasibility, we

require $0 < \rho < 1$, thus $\lambda \in (0, \ln(p/(1-p)))$. Now, since τ is a positive integer random variable, we can write:

$$\Rightarrow \mathbb{E}[\tau] = \sum_{n=0}^{\infty} \mathbb{P}[\tau > n] \le e^{\lambda}/(1-\rho)$$

The above bound is a function of the dummy variable λ , and the two MGFs f_1, f_2 . Our objective is to find λ and ρ that minimize this bound. We solve this numerically and compare the minimum bound with the mean time obtained by Monte-Carlo simulations for different values of p and δ .

B. Lower bound on $\mathbb{E}[\tau]$

Let $\tilde{\rho} = 1/\max(f_1(\lambda), f_2(\lambda))$ for regions where $0 < \tilde{\rho} < 1$, i.e. $\lambda \in (\ln(p/(1-p)), \infty)$. The following Proposition provides a lower bound on $\mathbb{E}[\tau]$.

Proposition 2. $\mathbb{E}[\tau] \geq e^{-\lambda} \tilde{\rho}[1-A]/[M_1(1-\tilde{\rho})]$, where $0 < \tilde{\rho} < 1$, $\lambda \in (\ln(p/(-p)), \infty)$,

$$A = \tilde{\rho}^4 + (\tilde{\rho} - \tilde{\rho}^4)\mathcal{P}_1 + (\tilde{\rho}^2 - \tilde{\rho}^4)\mathcal{P}_2 + (\tilde{\rho}^3 - \tilde{\rho}^4)\mathcal{P}_3, \text{ and}$$

$$\mathcal{P}_n = \mathbb{P}[\tau = n|h_1] \text{ for } n = 1, 2, 3.$$

Proof Idea. The proof uses the super-martingale property of $\{M_n\}$ and total probability theorem using the three possible values of h_1 . Conditioned on each h_1 , we calculate the probabilities of τ taking the first three values n=1,2,3. We then use these probabilities to provide a lower bound on $\mathbb{E}[\tau]$. See Appendix for details.

Note that the above lower bound is then numerically maximized over $\lambda \in (\ln(p/(1-p),\infty))$. Moreover, due to computational constraints, the bound in Proposition 2 is obtained using the closed-form expressions of \mathcal{P}_n for n=1,2,3. Next, we present an algorithm that improves this lower bound by numerically calculating \mathcal{P}_n for higher values of n.

Algorithm 2 Finding $\mathbb{P}[\tau = n|h_1]$

Input: $V = 0, p, \delta, n$ Output: $\mathcal{P}_n = \mathbb{P}[\tau = n|h_1]$

idea: Build a breadth-first tree conditioned on h_1 , tree.add(root), qualified = empty list of qualified nodes **while** tree.notempty() **do**

Pick the first node j at lowest level i by BFS Check for early elimination, e.g. $h_j > 1 + (n-i)/x$ if i < n (not a leaf) then check for $h_j \ge -1$ if True then tree.add(j's children) update condition on node j's children else(leaf) check for $h_j < -1$ if True then qualified.add(j)

tree.remove(j)

Update \mathcal{P}_n using the qualified list of leaves, return \mathcal{P}_n

C. Numerical and Simulation results

In Fig. 4 below, we use Algorithm 2 to show how the lower bound can be improved as n is increased. Conditioned on each h_1 , computational constraints limit us to using at

most n=17, which generates approximately 10^5 possible realizations of the history that would lead to an N cascade. The algorithm offers more improvement for lower values of δ . The non-monotonicities and discontinuities of the bound are a consequence of the same behaviors of each \mathcal{P}_n .

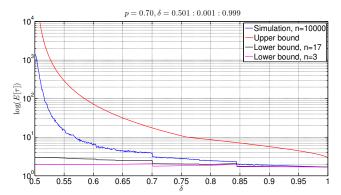


Figure 4: Bounds of $\log(\mathbb{E}[\tau])$ versus simulation.

Fig. 4 also shows the numerical bounds as compared with simulations on log-scale. Simulation results showed that $\mathbb{E}[\tau]$ is decreasing as δ increases. As $\delta \to 1$, the lower bound offers a better approximation to simulations; both converge to the same value of 1+p. On the other hand as $\delta \to 0.5$, both the upper bound and the simulation results blow up, while the lower bound offers less information.

In fact, we can verify that as $\delta \to 0.5$, we have $\mathbb{E}[\tau] \to \infty$. For any $\delta \in [0.5, 1]$, there is a positive probability of the underlying Markov chain hitting a state in the region where $h_n > 1$. In this region, the process becomes a simple birthdeath process with transition probabilities δ , $1 - \delta$ to the left and right, respectively. By relabeling the states, assume that we start at a state i > 0 in this birth-death process where 0 is an absorbing state on the left and there is no absorbing state on the right. Let $\tau_i = \min\{n > 0 : h_n = 0 | h_0 = i\}$ be the stopping time when the absorbing state is 0. For this birth-death process, the recurrence equation is written as: $(1-\delta)\mathbb{E}[au_{i+1}]-\mathbb{E}[au_i]+\delta\mathbb{E}[au_{i-1}]=-1.$ This generates a general solution of the form $\mathbb{E}[\tau_i] = A\left(\frac{\delta}{1-\delta}\right)^i + B + \frac{i}{2\delta-1}$, where A, B are real constants. Using the boundary condition $\mathbb{E}[\tau_0] = 0$, we have A = -B. Moreover, since $\mathbb{E}[\tau_i] > 0$, we require $A \ge 0$. Now assume that $\delta = 0.5 + \epsilon$ where we let $\epsilon \to 0$. As a result, $2\delta - 1 = 2\epsilon \to 0$ and $\mathbb{E}[\tau_i] \to \infty$. But since $\mathbb{E}[\tau_i]$ gives a lower bound on the original $\mathbb{E}[\tau]$, we also have $\mathbb{E}[\tau] \to \infty$ as $\delta \to 0.5$.

V. CONCLUSIONS AND FUTURE WORK

This paper studied a Bayesian learning model with information cascades. We assumed that subsequent agents can observe perfectly the previous actions and, in addition, feedback in the form of noisy reviews depending on the actions. We showed that the probabilities that agents cascade toward the wrong actions are not monotonic in the reviews quality. In particular, noisy reviews could increase the probability that agents misinterpret the true value of a good product. In practice, in online platforms like Yelp, Amazon, etc.

customers reviews come with a variability of strengths. Even though this scenario was not considered in this paper and our previous work, our results indirectly implied that a platform planner should opt to cut out the reviews of bad qualities and release only the truthful ones. In fact, this strategy is already adopted by those platforms, e.g. Amazon with verified purchase reviews, or Yelp with filtered reviews. Moreover, our results suggested that no matter how strong the reviews are improved to, agents might not perform better if their prior knowledge are limited. This implies that a platform planner should consider spending their budget on improving both the product's marketing efficiency and the reviews' reliability.

In the future work, we plan to study the possibility of having reviews with strengths non-homogeneously distributed across the population. In addition, we would like to study the probability of wrong cascades for more generic relationships between the signals quality and the reviews strength. Other possible directions include considering having reviews when both type of actions are taken, letting agents have the option to leave the reviews, and assuming that not all agents would exercise this option.

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APPENDIX

A. Proof sketch of Lemma 1

For simplicity of notation, the common denominator i of the accessible states is implicitly understood and we only consider the dynamics of the numerators. Therefore we reformulate $\mathscr{A}=\{-i,-i+1,\ldots,i-1,i\}$ with the initial state $h_0=0$. If $h_n\in[-i,i]$, the next state is h_n+X where $X\in\{-i,i+j,i-j\}$. If $h_n>i$, the next state is h_n+Y where $Y\in\{\pm-j\}$. If $h_n<-i$, h_n is absorbing where N cascade happens and the next state is also h_n .

1) For j > 3i, first verify after any number of steps, $\mathscr{B} \stackrel{\triangle}{=} \{-i, 0, i, 2i, i + aj, 2i + bj\}$, where integers $a, b \geq 1$, is the

set of all accessible states that are not absorbing. Moreover, from any state in \mathcal{B} , the next state is either another element of \mathcal{B} , or an absorbing state below -i. This can be showed using the fact that j > 3i.

For $j \in \{3i, 2i, i\}$, since gcd(i, j) = 1 one can simply pick i = 1 and j = 3, 2, 1 respectively. This proves that the accessible states in \mathscr{A} are limited to $\{-i, 0, i\}$.

2) For 3i > j > 2i, let j = 2i + z where 0 < z < i; and let i = kz + y where $0 \le y < z$ (note that y = 0 when z = 1). Let $\mathscr{C} \stackrel{\triangle}{=} \{-i, 0, i, -z, -2z, \dots, -kz, i - z, i - 2z, \dots, i - kz\}$.

The proof is done in 2 steps. First, verify that all states in $\mathscr C$ are accessible. This can be showed by considering the possible increments. Notice that from i one can access -z, and from -az one can access both i-az and -(a+1)z for integers $a=1,\ldots,k-1$. Finally, i-kz is accessible from -kz. Secondly, one can show that all accessible states are exclusively in $\mathscr C$. This is showed by exploring possible next states starting from any state $c\in\mathscr C$. Eventually, either an absorbing state is reached, or one ends up at another state in $\mathscr C$.

3) For the first case when 2i > j > i, let j = i + z where 0 < z < i. Since gcd(i,j) = 1, gcd(z,i) = 1 too. The idea is that, since gcd(z,i) = 1, the following i integers have different remainders when divided by i: $0, -z, -2z, \ldots, -(i-1)z$. Thus, if we can find a set of such i states that are accessible, and such that all those i states are in [0,i), then this implies that those i states are exactly $\{0,1,\ldots,i-1\}$. Moreover, it is obvious that state i accessible, thus all states in [0,i] are accessible. As a results, all states in [-i,0] are also accessible by adding an increment of -i.

For the second case when i>j, let i=kj+z where 0< z< j. Since gcd(i,j)=1, gcd(z,j)=1 too. The idea is similar to the first case, but using modulus j instead of modulus i. First, one needs to show that there exists an accessible set $\mathscr D$ of j states in [0,i] such that for $d_1,d_2\in\mathscr D$ and $d_1\ne d_2$, we have $d_1\not\equiv d_2\pmod j$. This says that the elements in $\mathscr D$, when divided by j, fill up all the possible remainders $0,1,\ldots,j-1$. Next, one needs to show there exists paths from those j states of $\mathscr D$ to j corresponding states in [0,j), which are exactly $0,1,\ldots,j-1$. Moreover, it is obvious that state j accessible, thus all states in [0,j] are accessible. Finally, one can always add/subtract integer multiples of j to access all states in [-i,i].

B. Proof of Proposition 3

Conditioning on h_1 , there are three possibilities: 1) $S_1=L$ gives $h_1=-1, M_1=\mathrm{e}^{-\lambda_2}\tilde{\rho};$ 2) $S_1=H, R_1=B$ gives $h_1=1-1/x, M_1=\mathrm{e}^{\lambda_2(1-1/x)}\tilde{\rho};$ and 3) $S_1=H, R_1=G$ gives $h_1=1+1/x, M_1=\mathrm{e}^{\lambda_2(1+1/x)}\tilde{\rho},$ and so we can write:

$$\mathbb{E}[\tau] = \sum_{h_1} \mathbb{P}[h_1] \mathbb{E}[\tau | h_1]$$
 (5)

Let $\tilde{\rho}=1/\max(f_1(\lambda),f_2(\lambda))$ for regions where $0<\tilde{\rho}<1$, i.e. $\lambda\in(\ln(\frac{p}{1-p}),\infty)$. Since τ is the first time $h_n<-1$,

it follows that $M_n \ge e^{-\lambda_2} \tilde{\rho}^n$ for $1 \le n \le \tau$, which yields:

$$\sum_{n=1}^{\tau} M_n \ge e^{-\lambda_2} \sum_{n=1}^{\tau} \tilde{\rho}^n = e^{-\lambda_2} \left[\frac{1 - \tilde{\rho}^{\tau+1}}{1 - \tilde{\rho}} - 1 \right]$$

Taking expectation on both sides (conditioned on h_1). and using the super-martingale property (so that $\mathbb{E}[M_n] \leq M_1$) we have:

$$\mathbb{E}[\tau|h_1] \ge e^{-\lambda_2} \frac{1}{M_1} \left[\frac{1 - \tilde{\rho} \mathbb{E}[\tilde{\rho}^{\tau}|h_1]}{1 - \tilde{\rho}} - 1 \right]$$
 (6)

We next find find an upper bound for $\mathbb{E}[\tilde{\rho}^{\tau}|h_1]$. Let $\mathcal{P}_n = \mathbb{P}[\tau = n|h_1]$ for $n = 1, 2, \dots$ We can rewrite:

$$\mathbb{E}[\tilde{\rho}^{\tau}|h_1] = \sum_{n=1}^{\infty} \mathcal{P}_n \tilde{\rho}^n = \mathcal{P}_1 \tilde{\rho} + \mathcal{P}_2 \tilde{\rho}^2 + \mathcal{P}_3 \tilde{\rho}^3 + \sum_{n=4}^{\infty} \mathcal{P}_n \tilde{\rho}^n,$$

where

$$\sum_{n=4}^{\infty} \mathcal{P}_n \tilde{\rho}^n \leq \tilde{\rho}^4 \sum_{n=4}^{\infty} \mathcal{P}_n = \tilde{\rho}^4 \left(1 - \mathcal{P}_1 - \mathcal{P}_2 - \mathcal{P}_3 \right).$$

Substitute this back to (6) and (5) we have the given bound in the Proposition.